

# Model Checking Stochastic Branching Processes

Taolue Chen   Klaus Dräger   *Stefan Kiefer*

University of Oxford, UK

MFCS 2012, Bratislava  
27 August 2012

Two classical model-checking problems:

- 1 Does a given **non-deterministic transition system** satisfy a given property?
- 2 What's the probability that a given **Markov chain** satisfies a given property?

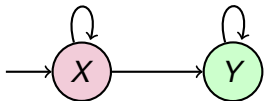
We consider linear-time properties:

$\omega$ -regular specifications, e.g., LTL formulae.

Our plan:

- Define a natural generalisation of those problems.
- Solve the generalised problem.

# Nondeterministic Transition Systems

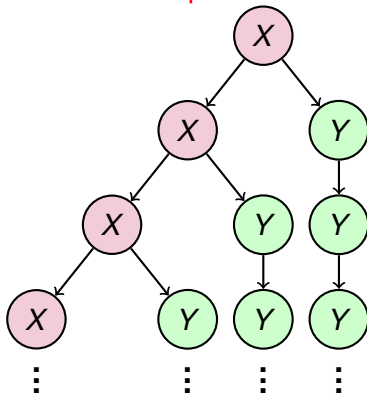


textual representation:

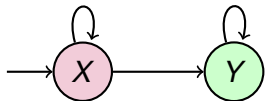
$$X \hookrightarrow XY \quad Y \hookrightarrow Y$$

(one rule for each state)

induces a **unique** tree:



# Nondeterministic Transition Systems

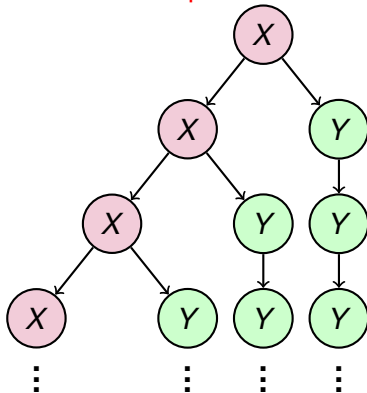


textual representation:

$$X \hookrightarrow XY \quad Y \hookrightarrow Y$$

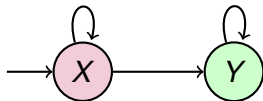
(one rule for each state)

induces a **unique** tree:



- Do all branches of the tree satisfy  $\Box(Y \rightarrow \Box Y)$ ?

# Nondeterministic Transition Systems

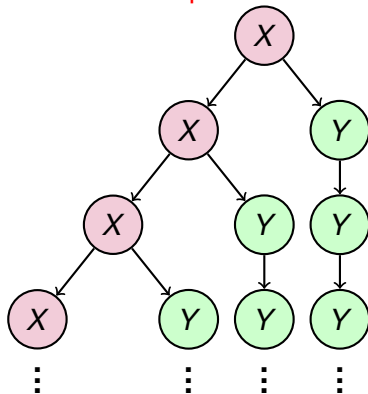


textual representation:

$$X \hookrightarrow XY \quad Y \hookrightarrow Y$$

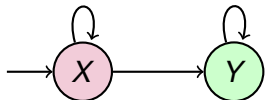
(one rule for each state)

induces a **unique** tree:



- Do all branches of the tree satisfy  $\Box(Y \rightarrow \Box Y)$ ? Yes.

# Nondeterministic Transition Systems

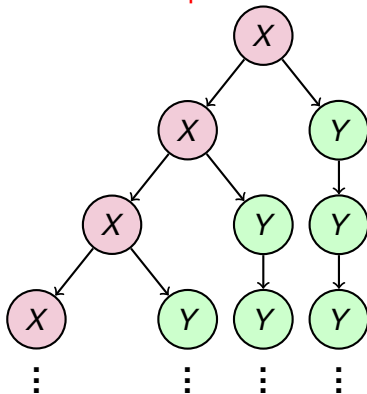


textual representation:

$$X \hookrightarrow XY \quad Y \hookrightarrow Y$$

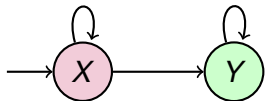
(one rule for each state)

induces a **unique** tree:



- Do all branches of the tree satisfy  $\Box(Y \rightarrow \Box Y)$ ? Yes.
- Do all branches of the tree satisfy  $\Diamond Y$ ?

# Nondeterministic Transition Systems

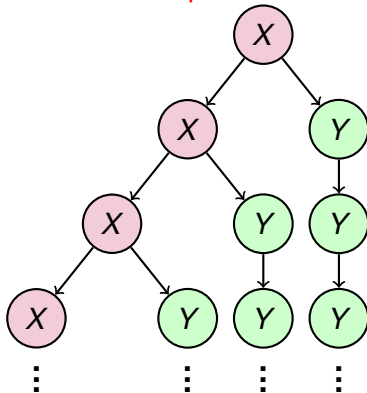


textual representation:

$$X \hookrightarrow XY \quad Y \hookrightarrow Y$$

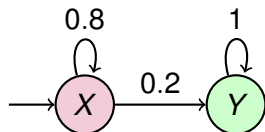
(one rule for each state)

induces a **unique** tree:



- Do all branches of the tree satisfy  $\Box(Y \rightarrow \Box Y)$ ? Yes.
- Do all branches of the tree satisfy  $\Diamond Y$ ? No.

# Markov Chains



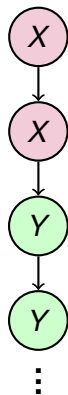
textual representation:

$$X \xrightarrow{0.8} X \quad Y \xrightarrow{1} Y$$

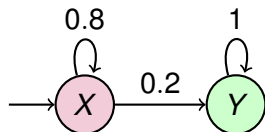
$$X \xrightarrow{0.2} Y$$

(multiple rules for each state)

induces a **random** “tree”  
(only one branch):







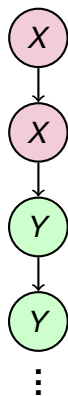
textual representation:

$$X \xrightarrow{0.8} X \quad Y \xrightarrow{1} Y$$

$$X \xrightarrow{0.2} Y$$

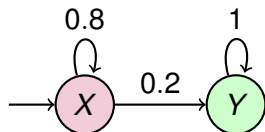
(multiple rules for each state)

induces a **random** “tree”  
(only one branch):



- Does the branch satisfy  $\varphi_1 := \Box(Y \rightarrow \Box Y)$ ?

# Markov Chains



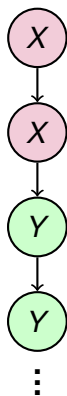
textual representation:

$$X \xrightarrow{0.8} X \quad Y \xrightarrow{1} Y$$

$$X \xrightarrow{0.2} Y$$

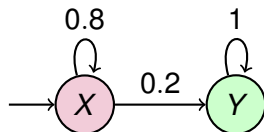
(multiple rules for each state)

induces a **random** “tree”  
(only one branch):



- Does the branch satisfy  $\varphi_1 := \Box(Y \rightarrow \Box Y)$ ?  $\Pr(\varphi_1) = 1$

# Markov Chains



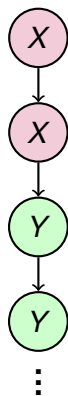
textual representation:

$$X \xrightarrow{0.8} X \quad Y \xrightarrow{1} Y$$

$$X \xrightarrow{0.2} Y$$

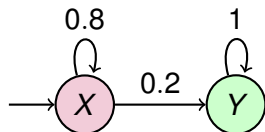
(multiple rules for each state)

induces a **random** “tree”  
(only one branch):



- Does the branch satisfy  $\varphi_1 := \Box(Y \rightarrow \Box Y)$ ?  $\Pr(\varphi_1) = 1$
- Does the branch satisfy  $\varphi_2 := \Diamond Y$ ?

# Markov Chains



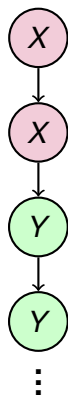
textual representation:

$$X \xrightarrow{0.8} X \quad Y \xrightarrow{1} Y$$

$$X \xrightarrow{0.2} Y$$

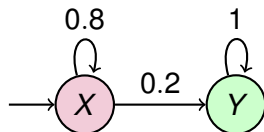
(multiple rules for each state)

induces a **random** “tree”  
(only one branch):



- Does the branch satisfy  $\varphi_1 := \Box(Y \rightarrow \Box Y)$ ?  $\Pr(\varphi_1) = 1$
- Does the branch satisfy  $\varphi_2 := \Diamond Y$ ?  $\Pr(\varphi_2) = 1$

# Markov Chains



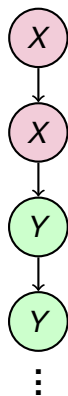
textual representation:

$$X \xrightarrow{0.8} X \quad Y \xrightarrow{1} Y$$

$$X \xrightarrow{0.2} Y$$

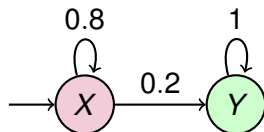
(multiple rules for each state)

induces a **random** “tree”  
(only one branch):



- Does the branch satisfy  $\varphi_1 := \Box(Y \rightarrow \Box Y)$ ?  $\Pr(\varphi_1) = 1$
- Does the branch satisfy  $\varphi_2 := \Diamond Y$ ?  $\Pr(\varphi_2) = 1$
- Does the branch satisfy  $\varphi_3 := \bigcirc Y$ ?

# Markov Chains



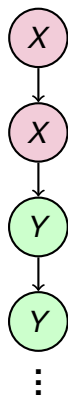
textual representation:

$$X \xrightarrow{0.8} X \quad Y \xrightarrow{1} Y$$

$$X \xrightarrow{0.2} Y$$

(multiple rules for each state)

induces a **random** “tree”  
(only one branch):



- Does the branch satisfy  $\varphi_1 := \Box(Y \rightarrow \Box Y)$ ?  $\Pr(\varphi_1) = 1$
- Does the branch satisfy  $\varphi_2 := \Diamond Y$ ?  $\Pr(\varphi_2) = 1$
- Does the branch satisfy  $\varphi_3 := \bigcirc Y$ ?  $\Pr(\varphi_3) = 0.2$

# Branching Processes

nondeterministic transition system:

degenerated probability distribution on trees  
(probability 1 for one tree, probability 0 for all others)

Markov chain:

probability distribution on degenerated trees  
(every node has just one child)

branching process:

probability distribution on trees

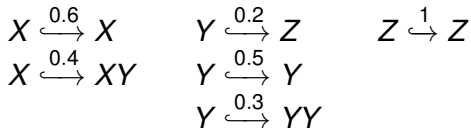
$$\begin{array}{l} X \xrightarrow{0.6} X \\ X \xrightarrow{0.4} XY \end{array}$$

$$\begin{array}{l} Y \xrightarrow{0.2} Z \\ Y \xrightarrow{0.5} Y \\ Y \xrightarrow{0.3} YY \end{array}$$

$$Z \xrightarrow{1} Z$$

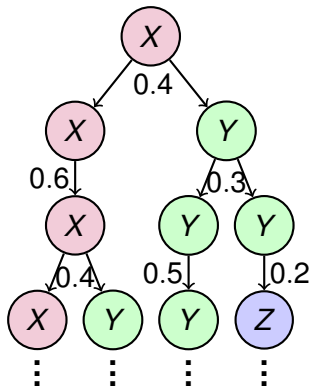
→ New plan: model check random (infinite) trees!

# Branching Processes



Probability of a tree that starts  
as on the right  
 $= 0.4 \cdot 0.6 \cdot 0.3 \cdot 0.4 \cdot 0.5 \cdot 0.2$

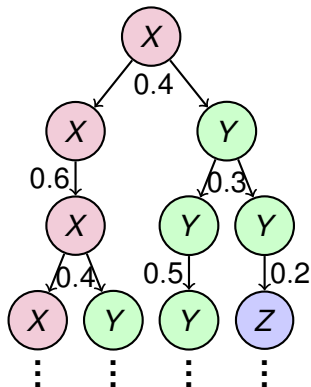
→ probability measure on (infinite) trees





# Model Checking: Simple Doesn't Work

$$\begin{array}{lll} X \xrightarrow{0.6} X & Y \xrightarrow{0.2} Z & Z \xrightarrow{1} Z \\ X \xrightarrow{0.4} XY & Y \xrightarrow{0.5} Y & \\ & Y \xrightarrow{0.3} YY & \end{array}$$

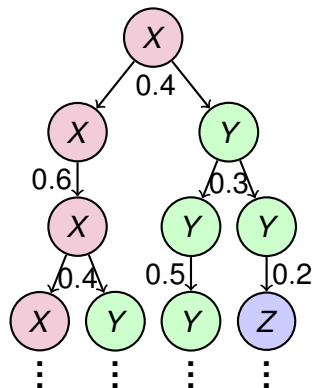


Consider  $\varphi := \square\lozenge(X \vee Z)$  (on all branches).  
What is  $\Pr_X(\varphi)$  ?

- “Markov-Chain” approach:  $\Pr_X(\varphi) = 1$
- “Pushdown-System” approach:  $\Pr_X(\varphi) = 1$

# Model Checking: Simple Doesn't Work

$$\begin{array}{lll} X \xrightarrow{0.6} X & Y \xrightarrow{0.2} Z & Z \xrightarrow{1} Z \\ X \xrightarrow{0.4} XY & Y \xrightarrow{0.5} Y & \\ & Y \xrightarrow{0.3} YY & \end{array}$$

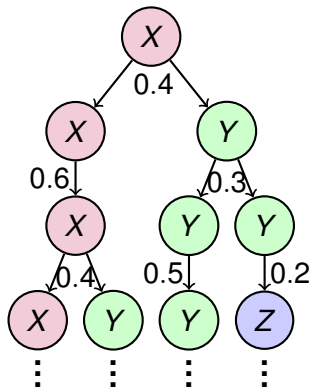


Consider  $\varphi := \square\Diamond(X \vee Z)$  (on all branches).  
What is  $\Pr_X(\varphi)$  ?

- “Markov-Chain” approach:  $\Pr_X(\varphi) = 1$
- “Pushdown-System” approach:  $\Pr_X(\varphi) = 1$
- correct value:  $\Pr_X(\varphi) = 0$

# Model Checking: Simple Doesn't Work

$$\begin{array}{lll} X \xrightarrow{0.6} X & Y \xrightarrow{0.2} Z & Z \xrightarrow{1} Z \\ X \xrightarrow{0.4} XY & Y \xrightarrow{0.5} Y & \\ & Y \xrightarrow{0.3} YY & \end{array}$$

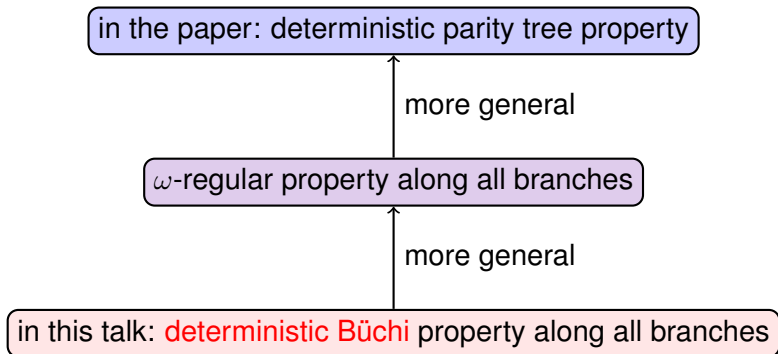


Consider  $\varphi := \Box\Diamond(X \vee Z)$  (on all branches).  
What is  $\Pr_X(\varphi)$  ?

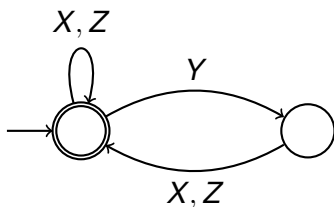
- “Markov-Chain” approach:  $\Pr_X(\varphi) = 1$
- “Pushdown-System” approach:  $\Pr_X(\varphi) = 1$
- correct value:  $\Pr_X(\varphi) = 0$

**However:** Swapping 0.2 and 0.3 in the rules  $\longrightarrow \Pr_X(\varphi) = 1$ .  
 $\Rightarrow$  The exact **numbers** matter, even for qualitative behaviour.

# Properties



# Deterministic Büchi Automata



- Perform a product construction with the branching process.  
(instance of automata-theoretic approach)
- Obtain a branching process with  
**accepting states** and **non-accepting states**.

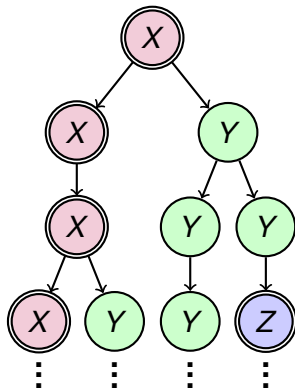
# Branching Process with (Non-)Accepting States

After product construction:

$$\begin{array}{l} X \xrightarrow{0.6} X \quad Y \xrightarrow{0.2} Z \quad Z \xrightarrow{1} Z \\ X \xrightarrow{0.4} XY \quad Y \xrightarrow{0.5} Y \\ \quad \quad \quad Y \xrightarrow{0.3} YY \end{array}$$

Accepting: X, Z

Non-Accepting: Y



a tree is **good**  $\stackrel{\text{def}}{\iff}$

each branch has infinitely many accepting nodes

Compute  $\Pr_X(\text{good})$

# Decent Trees

a tree is **good**  $\stackrel{\text{def}}{\iff}$   
each branch has  $\infty$  accepting nodes

a tree is **decent**  $\stackrel{\text{def}}{\iff}$   
each branch has  $\geq 1$  accepting node (besides the root)

$$X \xrightarrow{0.8} YY$$

$$X \xrightarrow{0.2} Z$$

X non-acc.

$$Y \xrightarrow{0.3} XX$$

$$Y \xrightarrow{0.7} Z$$

Y non-acc.

$$Z \xrightarrow{1} X$$

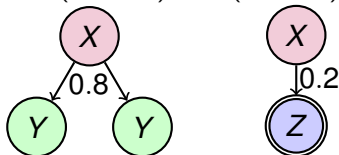
Z acc.

Equation system for  $\Pr_X(\text{decent})$ ,  $\Pr_Y(\text{decent})$ ,  $\Pr_Z(\text{decent})$ :

$$x = 0.8y^2 + 0.2$$

$$y = 0.3x^2 + 0.7$$

$$z = x$$



# Decent Trees

a tree is **good**  $\stackrel{\text{def}}{\iff}$   
each branch has  $\infty$  accepting nodes

a tree is **decent**  $\stackrel{\text{def}}{\iff}$   
each branch has  $\geq 1$  accepting node (besides the root)

$$X \xrightarrow{0.8} YY$$

$$Y \xrightarrow{0.3} XX$$

$$Z \xrightarrow{1} X$$

$$X \xrightarrow{0.2} Z$$

$$Y \xrightarrow{0.7} Z$$

X non-acc.

Y non-acc.

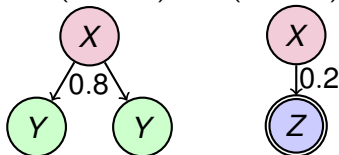
Z acc.

Equation system for  $\Pr_X(\text{decent})$ ,  $\Pr_Y(\text{decent})$ ,  $\Pr_Z(\text{decent})$ :

$$x = 0.8y^2 + 0.2$$

$$y = 0.3x^2 + 0.7$$

$$z = x$$



The **least solution** gives the correct probabilities.



# Solving Polynomial Fixed-Point Equations

$$\begin{aligned}x &= 0.8y^2 + 0.2 \\ y &= 0.3x^2 + 0.7\end{aligned}$$

One can “efficiently compute” the least solution:

Theorem (from literature, especially [ESY, STOC'12])

Let  $\vec{x} = \vec{f}(\vec{x})$  be an equation system where

- $\vec{x}$  is a vector of variables
- $\vec{f}$  is a vector of polynomials with nonnegative coefficients
- $\vec{f}(\vec{1}) = \vec{1}$  ( $\vec{1}$  = vector of 1s)

Let  $q$  be the first entry of the least solution. Then:

- (a) One can decide in polynomial time: Is  $q = 0$ ? Is  $q = 1$ ?
- (b) One can decide in polynomial space: Is  $q \bowtie \tau$ ?  
( $\tau \in \mathbb{Q}$  and  $\bowtie \in \{<, >, \leq, \geq, =, \neq\}$ ).
- (c) One can approximate  $q$  within additive error  $2^{-j}$  in time polynomial in  $j$  and the representation size of  $\vec{f}$ .

## Corollary

*One can in polynomial time:*

- (a) *decide whether  $\Pr_X(\text{decent}) = 1$*
- (b) *“efficiently approximate”  $\Pr_X(\text{decent})$ .*

Proof. Set up the equation system and apply the theorem. □

## Corollary

*One can in polynomial time:*

- (a) *decide whether  $\Pr_X(\text{decent}) = 1$*
- (b) *“efficiently approximate”  $\Pr_X(\text{decent})$ .*

Proof. Set up the equation system and apply the theorem. □

We want to do the same for  $\Pr_X(\text{good})$ .

First focus on the qualitative problem: Is  $\Pr_X(\text{good}) = 1$  ?

# The Qualitative Problem

Generalise the notion of “decent”:

a tree is  **$k$ -decent**  $\stackrel{\text{def}}{\iff}$   
each branch has  $\geq k$  accepting nodes (besides the root)

Note:  $1$ -decent  $\equiv$  decent and  $\infty$ -decent  $\equiv$  good

a state  $X$  is  **$k$ -safe**  $\stackrel{\text{def}}{\iff} \Pr_X(k\text{-decent}) = 1$

One can compute the **1**-safe states:

- remember:  $X$  is 1-safe means  $\Pr_X(\text{decent}) = 1$
- apply the corollary

One can compute the **2**-safe states:

- modify the branching process:  
accepting  $:=$  accepting  $\wedge$  **1**-safe
- compute the **1**-safe states in that process

# The Qualitative Problem

Generalise the notion of “decent”:

a tree is  **$k$ -decent**  $\stackrel{\text{def}}{\iff}$   
each branch has  $\geq k$  accepting nodes (besides the root)

Note:  $1$ -decent  $\equiv$  decent and  $\infty$ -decent  $\equiv$  good

a state  $X$  is  **$k$ -safe**  $\stackrel{\text{def}}{\iff} \Pr_X(k\text{-decent}) = 1$

One can compute the  **$1$ -safe** states:

- remember:  $X$  is  $1$ -safe means  $\Pr_X(\text{decent}) = 1$
- apply the corollary

One can compute the  **$k$ -safe** states:

- modify the branching process:  
accepting  $:=$  accepting  $\wedge$   **$(k-1)$ -safe**
- compute the  **$1$ -safe** states in that process

# The Qualitative Problem

1-safe  $\supseteq$  2-safe  $\supseteq$  3-safe  $\supseteq$  ...

The sequence must stabilise after  $n$  iterations.

Hence:

state  $X$  is  $n$ -safe  $\iff X$  is  $\infty$ -safe  $\iff \Pr_X(\text{good}) = 1$

## Theorem (Qualitative Problem)

*One can decide in polynomial time whether  $\Pr_X(\text{good}) = 1$ .*

# The Quantitative Problem

## Theorem (Quantitative Problem)

*One can “efficiently approximate”  $\Pr_X(\text{good})$ .*

Proof sketch. The following algorithm works:

1. compute the  $\infty$ -safe states in polynomial time (that is the qualitative problem).
2. modify the branching process:  
accepting  $:= \infty$ -safe
3. approximate  $\Pr_X(\text{decent})$  for the resulting process

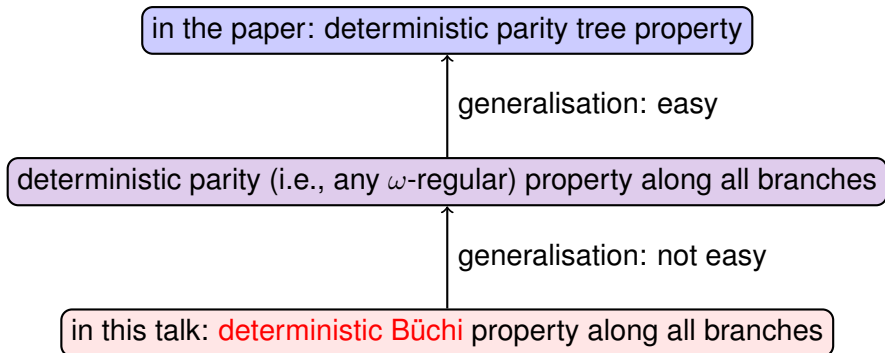
Correctness is not obvious, but not hard. □

deterministic parity (i.e., any  $\omega$ -regular) property along all branches

↑  
generalisation: not easy

in this talk: **deterministic Büchi** property along all branches





For **Markov Chains**: PCTL formulae

e.g.  $[\varphi U \psi]_{\geq 0.9} \equiv \Pr(\text{run satisfies } \varphi U \psi) \geq 0.9$

For **Branching Processes**: PTTL formulae

e.g.  $[\varphi EU \psi]_{\geq 0.9} \equiv \Pr(\text{tree has branch satisfying } \varphi U \psi) \geq 0.9$

e.g.  $[\varphi AU \psi]_{\geq 0.9} \equiv \Pr(\text{all branches satisfy } \varphi U \psi) \geq 0.9$

For **Markov Chains**: PCTL formulae

e.g.  $[\varphi U \psi]_{\geq 0.9} \equiv \Pr(\text{run satisfies } \varphi U \psi) \geq 0.9$

For **Branching Processes**: PTTL formulae

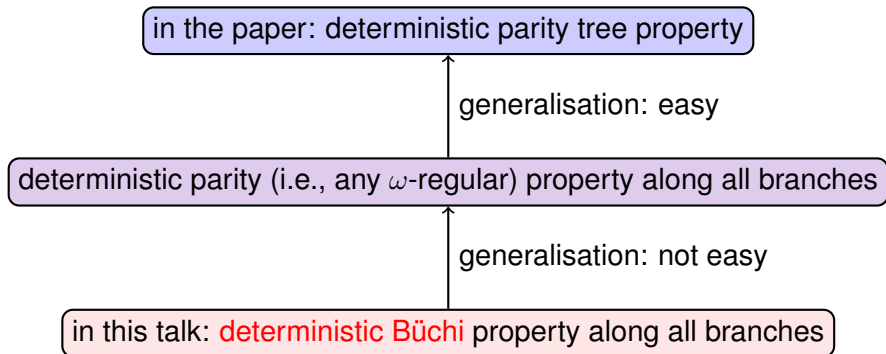
e.g.  $[\varphi EU \psi]_{\geq 0.9} \equiv \Pr(\text{tree has branch satisfying } \varphi U \psi) \geq 0.9$

e.g.  $[\varphi AU \psi]_{\geq 0.9} \equiv \Pr(\text{all branches satisfy } \varphi U \psi) \geq 0.9$

## Theorem

*Model checking branching processes*

- *against PTTL is in PSPACE.*
- *against the qualitative fragment of PTTL is in P.*



# Properties

